

Text book: T. Basar and G. Olsder "Dynamic Noncooperative Games," SIAM, 1999

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 on games

NORMAL GAME FORM

The extensive game form described by a game tree displays explicitly the evolution of the game. However, the description by itself becomes complex in the case of many players, many alternatives, and many stages.

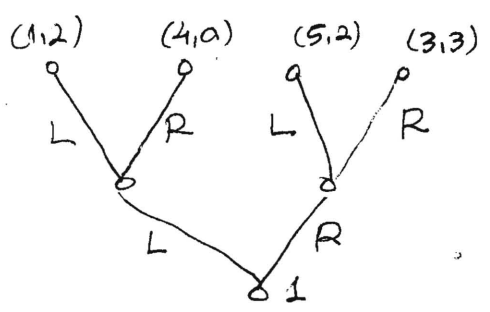
The normal game form suppresses information about game evolution and expresses the utility (loss) functions explicitly in terms of players' strategies.

$$J_i = J_i(u_1, u_2, \dots, u_N), \quad i = 1, 2, \dots, N$$

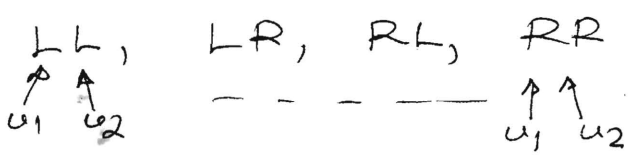
J_i = loss function (utility function, performance criterion)
 (cost function)

u_i = strategy of player i

(Example)



All possible strategies are



Hence

$$J_1 = J_1(u_1, u_2)$$

$$J_2 = J_2(u_1, u_2)$$

Rational behavior of each player is to minimize his (her) own losses assuming that the other player is doing the same.

$$J_1(u_1^*, u_2^*) \leq J_1(u_1, u_2^*)$$

$$J_2(u_1^*, u_2^*) \leq J_2(u_1^*, u_2)$$

Such a strategy (u_1^*, u_2^*) , assuming that it exists, is called the equilibrium strategy.

In general, for N -players, the loss functions are defined by

$$J_i = J_i(u_1, u_2, \dots, u_N) \quad , \quad i = 1, 2, \dots, N$$

with the equilibrium strategy satisfying

$$J_i(u_1^*, u_2^*, \dots, u_i^*, \dots, u_N^*) \leq J_i(u_1^*, u_2^*, \dots, u_i, u_{i+1}^*, \dots, u_N^*)$$

For $N=2$, normal form games are called the **MATRIX** games since they can be described using matrices

		player 2	
		L	R
player 1	L	(1, 2)	(4, 0)
	R	(5, 2)	(3, 3)

or

		P2	
		L	R
P1	L	$J_1=1$ $J_2=2$	$J_1=4$ $J_2=0$
	R	$J_1=5$ $J_2=0$	$J_1=3$ $J_2=3$

GAMES IN EUCLIDEAN SPACES (Infinite games)

1) ZERO-SUM GAMES

- a) simultaneous decision making \Rightarrow saddle point solution
- b) sequential decision making \Rightarrow minmax or maximum solution

2) CONFLICT GAMES

- a) simultaneous decision making \Rightarrow Nash strategies
- b) sequential decision making \Rightarrow Stackelberg "

3) COOPERATIVE GAMES \Rightarrow Pareto strategies

All games on E^m spaces are described by a static state equation

$$f(x, u) = 0$$

$x \in \mathbb{R}^m$ = state of the game vector

$u \in \mathbb{R}^m$ with $u = (u_1, u_2, \dots, u_M)$

$u_i \in \mathbb{R}^{m_i}$ = strategy of player i

In addition, each player has a loss function (utility function, cost, performance criterion)

$$J_i(x, u) \quad , \quad i = 1, 2, \dots, M$$

Assuming that $f(x, u) = 0 \Rightarrow x = \varphi(u)$

$$\Rightarrow J_i(\varphi(u), u) = J_i(u_1, u_2, \dots, u_M)$$

Note that none of the players has complete control over the game outcome.

1) ZERO-SUM GAMES ($J_2 = -J_1$)

a) Saddle point strategies

From

$$\left. \begin{aligned} J_1(u_1^*, u_2^*) &\leq J_1(u_1, u_2^*) \\ J_2(u_1^*, u_2^*) &\leq J_2(u_1^*, u_2) \end{aligned} \right\} (1)$$

and the fact that $J_2 = -J_1$, and

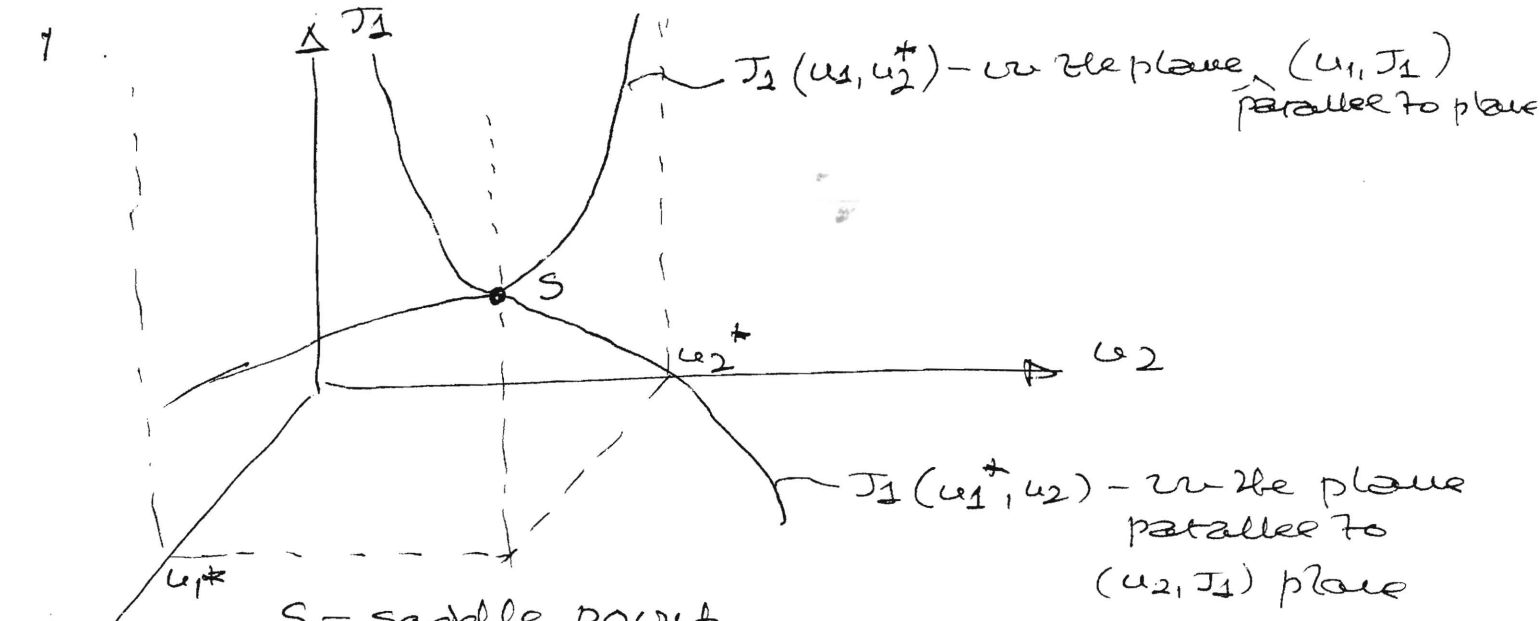
$$\begin{aligned} J_2(u_1^*, u_2^*) \leq J_2(u_1^*, u_2) &\implies -J_1(u_1^*, u_2^*) \leq -J_1(u_1^*, u_2) \\ &\implies J_1(u_1^*, u_2^*) \geq J_1(u_1^*, u_2) \end{aligned}$$

we get from (1)

$$\boxed{J_1(u_1^*, u_2) \leq J_1(u_1^*, u_2^*) \leq J_1(u_1, u_2^*)} \quad (2)$$

which represents the saddle point condition

Apparently in this zero-sum game the player 1 is minimizer and the player 2 is the maximizer.



S = saddle point

It exists if $J_1(u_1, u_2)$ is convex with respect to u_1 and concave with respect to u_2